

Number Representation and Arithmetic Operations

1.4.1 Integers

Consider an n -bit vector : $B = b_{n-1} \dots b_1 b_0$ where $b_i = 0$ or 1 for $0 \leq i \leq n-1$.

$$V(B) = b_{n-1} \times 2^{n-1} + \dots + b_1 \times 2^1 + b_0 \times 2^0$$

We need to represent both positive and negative numbers. Three systems are used for representing such numbers:

- Sign-and-magnitude
- 1's-complement
- 2's-complement

In all three systems, the leftmost bit is 0 for positive numbers and 1 for negative numbers.

B $b_3 b_2 b_1 b_0$	Values represented		
	Sign and magnitude	1's complement	2's complement
0 1 1 1	+7	+7	+7
0 1 1 0	+6	+6	+6
0 1 0 1	+5	+5	+5
0 1 0 0	+4	+4	+4
0 0 1 1	+3	+3	+3
0 0 1 0	+2	+2	+2
0 0 0 1	+1	+1	+1
0 0 0 0	+0	+0	+0
1 0 0 0	-0	-7	-8
1 0 0 1	-1	-6	-7
1 0 1 0	-2	-5	-6
1 0 1 1	-3	-4	-5
1 1 0 0	-4	-3	-4
1 1 0 1	-5	-2	-3
1 1 1 0	-6	-1	-2
1 1 1 1	-7	-0	-1

In *1's-complement* representation, negative values are obtained by complementing each bit of the corresponding positive number. Thus, the representation for -3 is obtained by complementing each bit in the vector 0011 to yield 1100.

In the *2's-complement* system, forming the 2's-complement of an n -bit number is done by subtracting the number from 2^n . Hence, the 2's-complement of a number is obtained by adding 1 to the 1's-complement of that number.

There are distinct representations for $+0$ and -0 in both the sign-and magnitude and 1's-complement systems, but the 2's-complement system has only one representation for 0.

Addition of Unsigned Integers

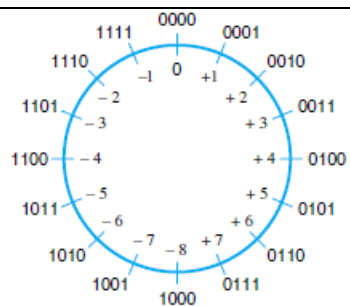
The sum of 1 and 1 is the 2-bit vector 10, which represents the value 2. We say that the *sum* is 0 and the *carry-out* is 1. We add bit pairs starting from the low-order (right) end of the bit vectors, propagating carries toward the high-order (left) end. The carry-out from a bit pair becomes the *carry-in* to the next bit pair to the left.

$$\begin{array}{cccc}
 0 & 1 & 0 & 1 \\
 + 0 & + 0 & + 1 & + 1 \\
 \hline
 0 & 1 & 1 & 10 \\
 & & & \uparrow \\
 & & & \text{Carry-out}
 \end{array}$$

Addition and Subtraction of Signed Integers

The 2's-complement system is the most efficient method for performing addition and subtraction operations.

Unsigned integers mod N is a circle with the values 0 through $N-1$. The decimal values 0 through 15 are represented by their 4-bit binary values 0000 through 1111.



(b) Mod 16 system for 2's-complement numbers

The operation $(7 + 5) \bmod 16$ yields the value 12. To perform this operation graphically, locate 7 (0111) on the outside of the circle and then move 5 units in the clockwise direction to arrive at the answer 12 (1100).

Similarly, $(9 + 14) \bmod 16 = 7$; this is modeled on the circle by locating 9 (1001) and moving 14 units in the clockwise direction past the zero position to arrive at the answer 7 (0111).

Apply the mod 16 addition technique to the example of adding +7 to -3. The 2's-complement representation for these numbers is 0111 and 1101, respectively.

To *add* two numbers, add their n -bit representations, ignoring the carry-out bit from the most significant bit (MSB) position. The sum will be the algebraically correct value in 2's-complement representation if the actual result is in the range -2^{n-1} through $+2^{n-1} - 1$.

To *subtract* two numbers X and Y , that is, to perform $X - Y$, form the 2's-complement of Y , then add it to X using the *add* rule. Again, the result will be the algebraically correct value in 2's-complement representation if the actual result is in the range -2^{n-1} through $+2^{n-1} - 1$.

$$\begin{array}{r} 0010 \quad (+2) \\ + 0011 \quad (+3) \\ \hline 0101 \quad (+5) \\ \\ 1011 \quad (-5) \\ + 1110 \quad (-2) \\ \hline 1001 \quad (-7) \end{array}$$

$$\begin{array}{r} (b) \quad 0100 \quad (+4) \\ + 1010 \quad (-6) \\ \hline 1110 \quad (-2) \\ \\ (d) \quad 0111 \quad (+7) \\ + 1101 \quad (-3) \\ \hline 0100 \quad (+4) \end{array}$$

Floating-Point Numbers

If we use a full word in a 32-bit word length computer to represent a signed integer in 2's-complement representation, the range of values that can be represented is -2^{31} to $+2^{31} - 1$.

Since the position of the binary point in a floating-point number varies, it must be indicated explicitly in the representation. For example, in the familiar decimal scientific notation, numbers may be written as 6.0247×10^{23} , 3.7291×10^{-27} , -1.0341×10^2 , -7.3000×10^{-14} . these numbers have been given to 5 *significant digits* of precision.

A binary floating-point number can be represented by:

- a sign for the number
- some significant bits
- a signed scale factor exponent for an implied base of 2

Character Representation

Bit positions	Bit positions 654							
3210	000	001	010	011	100	101	110	111
0000	NUL	DLE	SPACE	0	@	P	^	p
0001	SOH	DC1	!	1	A	Q	a	q
0010	STX	DC2	"	2	B	R	b	r
0011	ETX	DC3	#	3	C	S	c	s
0100	EOT	DC4	\$	4	D	T	d	t
0101	ENQ	NAK	%	5	E	U	e	u
0110	ACK	SYN	&	6	F	V	f	v
0111	BEL	ETB	'	7	G	W	g	w
1000	BS	CAN	(8	H	X	h	x
1001	HT	EM)	9	I	Y	i	y
1010	LF	SUB	*	:	J	Z	j	z
1011	VT	ESC	+	;	K	[k	{
1100	FF	FS	,	<	L	/	l	
1101	CR	GS	-	=	M]	m	}
1110	SO	RS	.	>	N	^	n	~
1111	SI	US	/	?	O	_	o	DEL

The most common encoding scheme for characters is ASCII (American Standard Code for Information Interchange). Alphanumeric characters, operators, punctuation symbols, and control characters are represented by 7-bit codes. It is convenient to use an 8-bit *byte* to represent and store a character.

The code occupies the low-order seven bits. The high-order bit is usually set to 0. This facilitates sorting operations on alphabetic and numeric data.

The low-order four bits of the ASCII codes for the decimal digits 0 to 9 are the first ten values of the binary number system.

This 4-bit encoding is referred to as the *binary-coded decimal* (BCD) code.